

COMPUTER SYSTEMS AND ORGANIZATION

Muxes

Daniel Graham



ENGINEERING

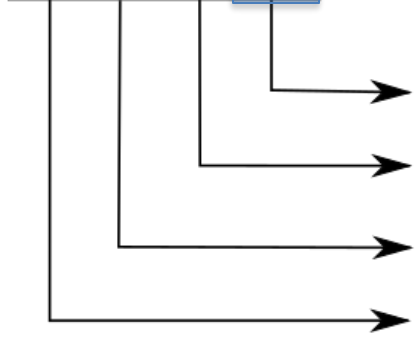
REVIEW

ENDIANNESS

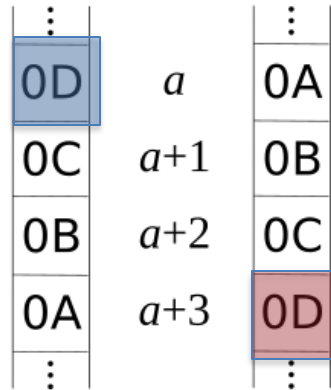
Little-endian

32-bit integer

0A0B0C0D



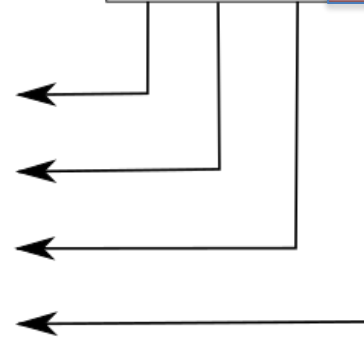
Memory



Big-endian

32-bit integer

0A0B0C0D



EXAM QUESTION FALL 2018

Question 13 [2 pt]: If the 32-bit number `0x12345678` is stored in **little-endian** at address `0x20`, what is the value of the byte at address `0x22`? Answer in hexadecimal.

Answer:

Question 14 [2 pt]: If you read the bytes `[fe, dc]` as an unsigned **big-endian** 16-bit number, what is that number? Answer in hexadecimal.

Answer:

EXAM QUESTION FALL 2018

Question 13 [2 pt]: If the 32-bit number `0x12345678` is stored in **little-endian** at address `0x20`, what is the value of the byte at address `0x22`? Answer in hexadecimal.

Answer:

`0x34`

Question 14 [2 pt]: If you read the bytes [`fe`, `dc`] as an unsigned **big-endian** 16-bit number, what is that number? Answer in hexadecimal.

Answer:

`0xfedc`

EXAM QUESTIONS FALL 2019

Suppose an array of two 32-bit values ($[0xabcdef01, 0x7645231]$) is stored at address $0x200$. What byte is stored at address $0x204$? Answer in hexadecimal.

Question 13 [2 pt]: (see above) Assume little-endian storage.

Answer:

Question 14 [1 pt]: (see above) Assume big-endian storage.

Answer:

Remember to draw a picture of how arrays are stored in memory.

(Remember to group second number into bytes)

EXAM QUESTIONS FALL 2019

Suppose an array of two 32-bit values ($[0xabcdef01, 0x7645231]$) is stored at address $0x200$. What byte is stored at address $0x204$? Answer in hexadecimal.

Question 13 [2 pt]: (see above) Assume little-endian storage.

Answer:

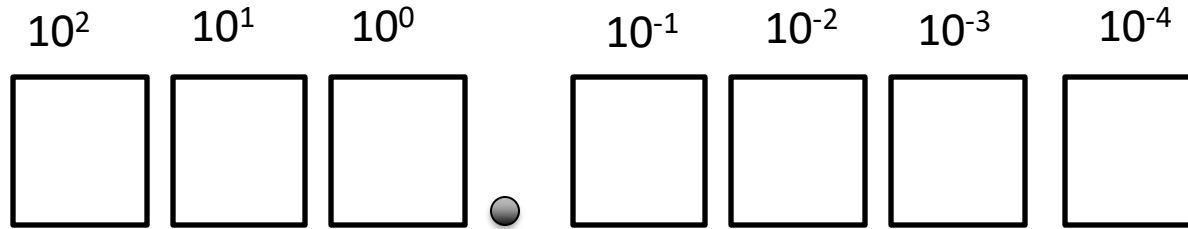
0x31

Question 14 [1 pt]: (see above) Assume big-endian storage.

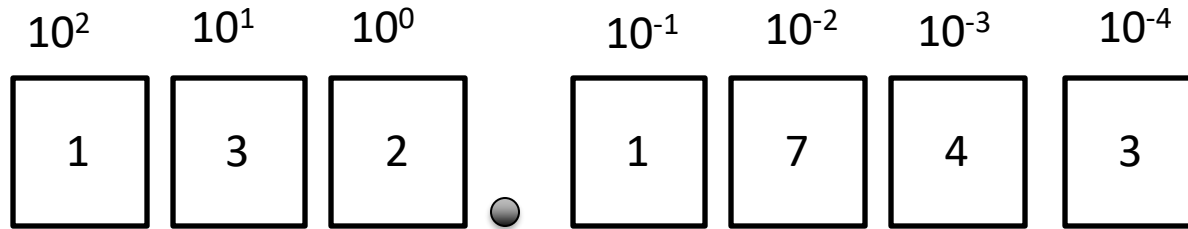
Answer:

0x07

FLOATING POINT BASE 10

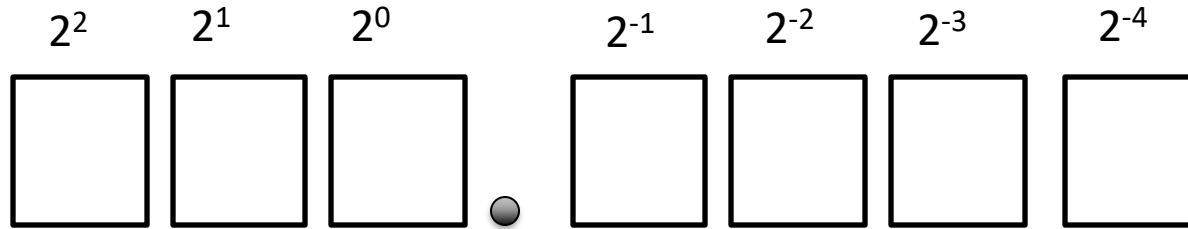


FLOATING POINT BASE 10

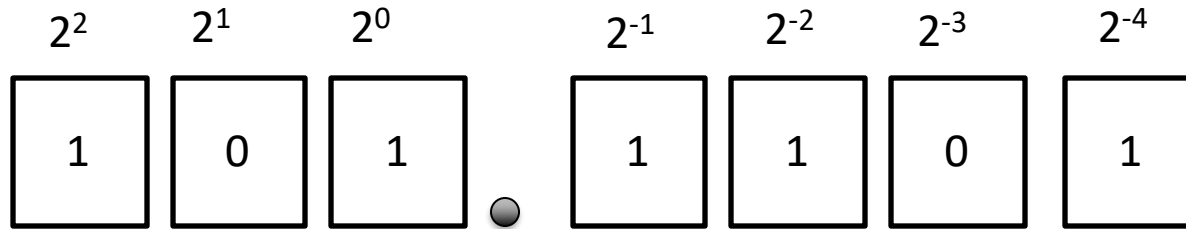


$$1 \times 10^2 + 3 \times 10^1 + 2 \times 10^0 + 1 \times 1/10 + 7 \times 1/100 + 4 \times 1/1000 + 3 \times 1/10000 = 132.1743$$

FLOATING POINT BASE 2



FLOATING POINT BASE 2



$$1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 5$$

$$1 \times 1/2 + 1 \times 1/4 + 0 \times 1/8 + 1 \times 1/16 = 13/16$$

$$5 \frac{13}{16} = 5 + 0.8125 = 5.8125$$

NOW WE JUST NEED THE EXPONENT



Now we need the exponent so we can float the point.

$$74.0 * 10^2$$

$$740.0 * 10^1$$

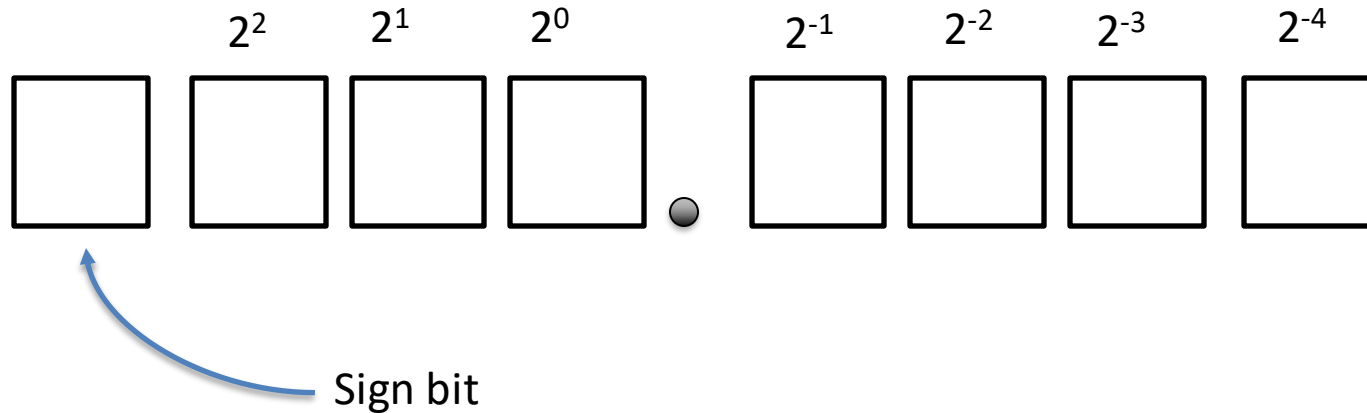
$$7400.0 * 10^1$$

$$-74.0 * 10^2$$

$$- 740.0 * 10^1$$

$$- 7400.0 * 10^1$$

FLOATING POINT BASE 2 (NEGATIVE NUMBERS)



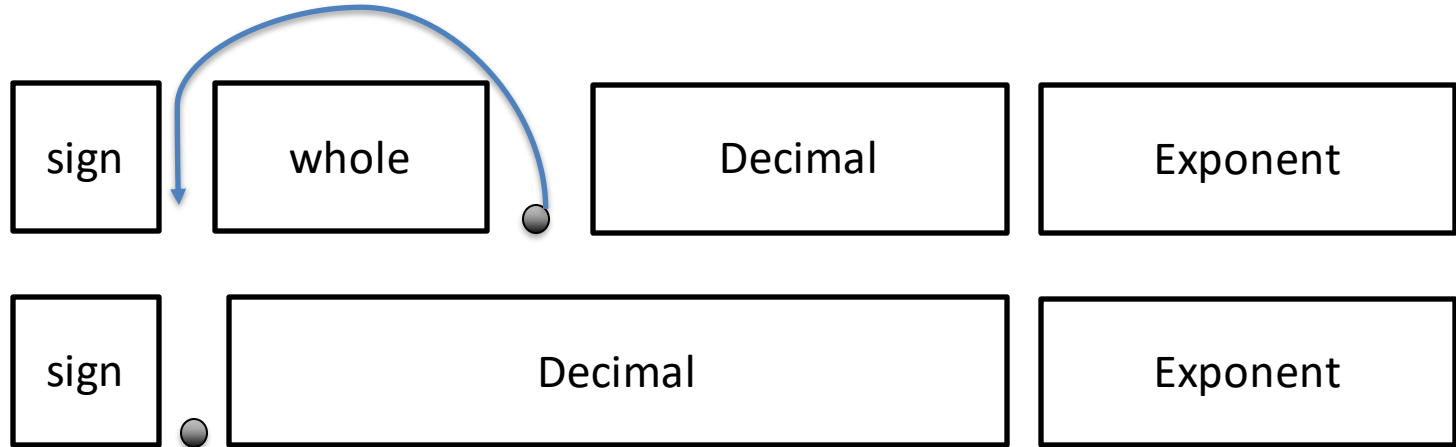
NOW WE JUST NEED THE EXPONENT



$$\text{number} = \text{sign}(\text{whole} + \text{decimal}) \times 2^{\text{exponent}}$$

$$1 \text{ 11.111} \times 2^E$$

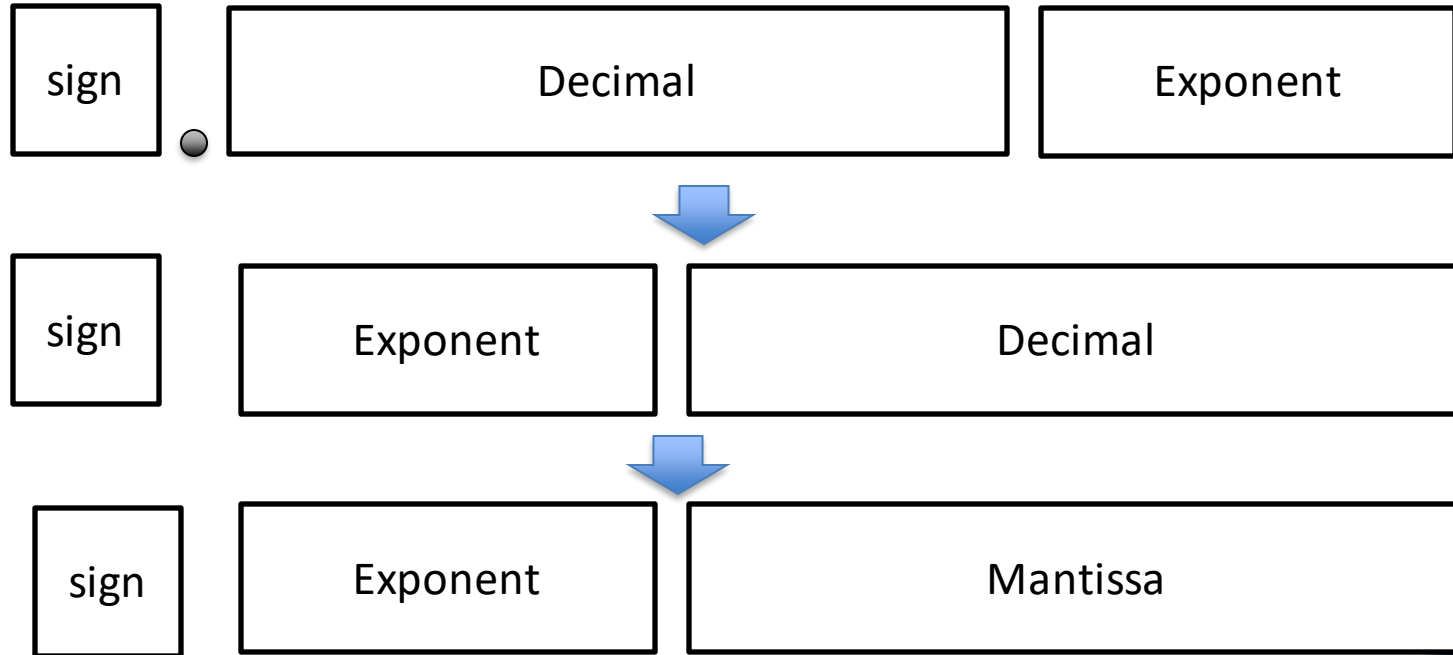
ADDING THE EXPONENT DELETING THE WHOLE NUMBER SECTION



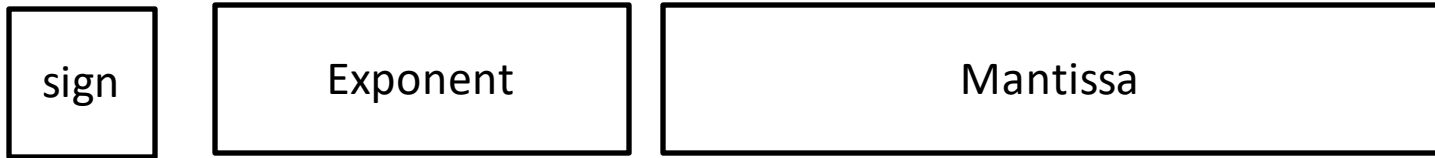
$$-74.01 * 10^2$$

$$- 0.7401 * 10^4$$

IEEE 754 FLOATING POINT STANDARD



IEEE 754

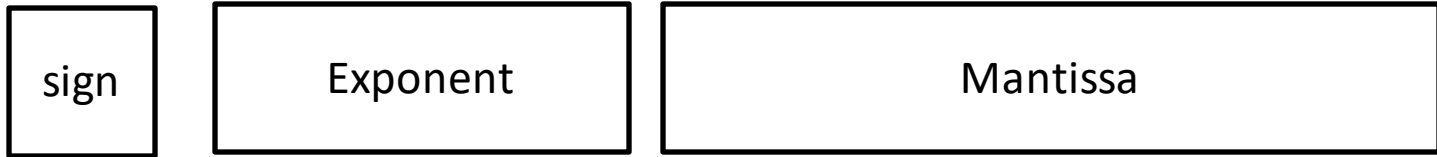


$$\text{number} = \text{sign}(1 + \text{Mantissa}) \times 2^{\text{exponent} - \text{bias}}$$

On 32-bit machines bias is normally 127 (Yes this is bias representation we talked about earlier)
0b01111111 (8 digits) is what we use when adding to the exponent

For 64-bit machines, the bias would be 1023, and this format allows for handling of negative exponents as the biased exponent is always stored as a positive number

IEEE 754



$$\text{number} = \text{sign}(1 + \text{Mantissa}) \times 2^{\text{exponent} - \text{bias}}$$

Remember this is a
base 2 binary string

$$1.\text{ffff} \times 2^{\text{exponent} - \text{bias}}$$

Remember 2^0 is
one

BINARY STRING

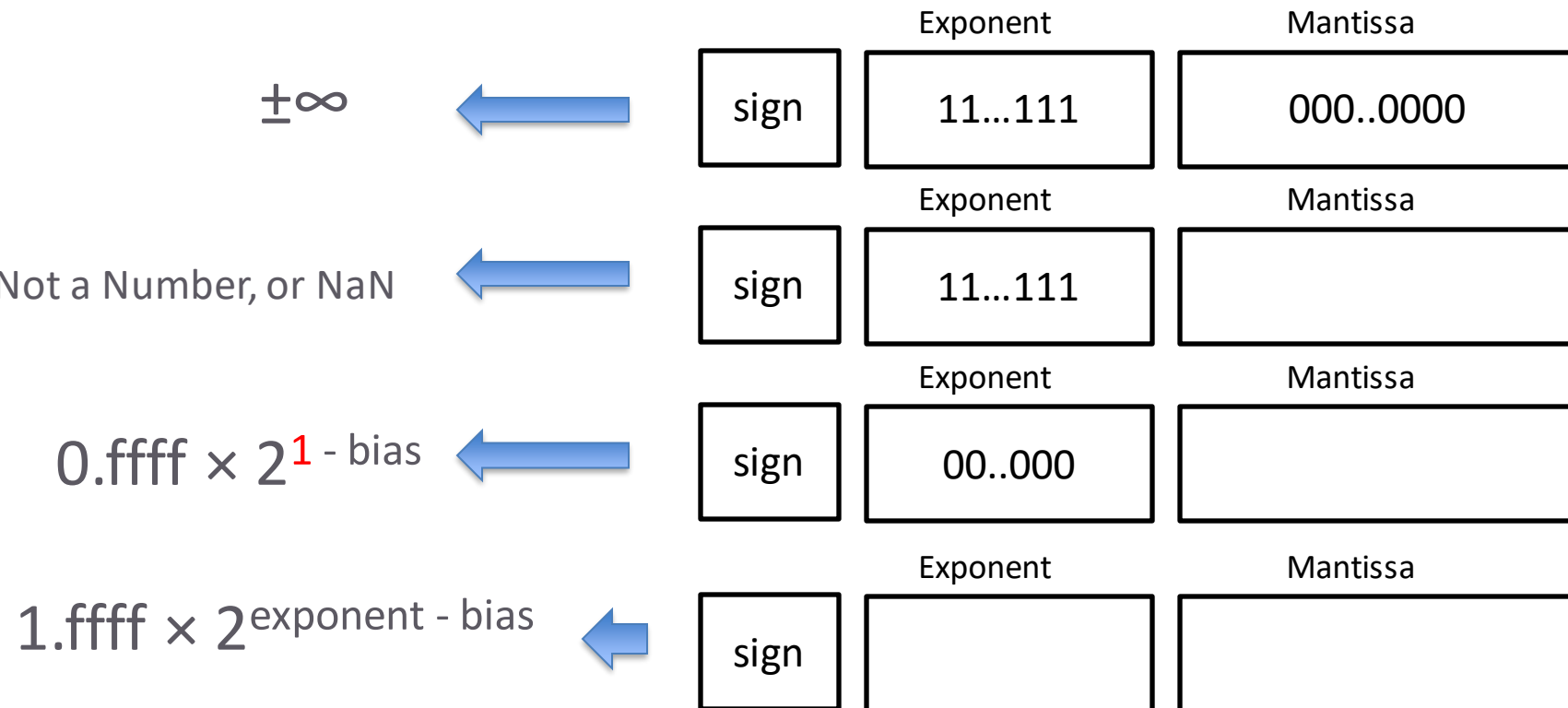
$$0.1101 = 1.101 \times 2^{-1}$$

$$0.01101 = 1.101 \times 2^{-2}$$

$$0.001101 = 1.101 \times 2^{-3}$$


Keep going until you get to your first 1. (one 1 left to the decimal point)

IEEE 754 FOUR CASES

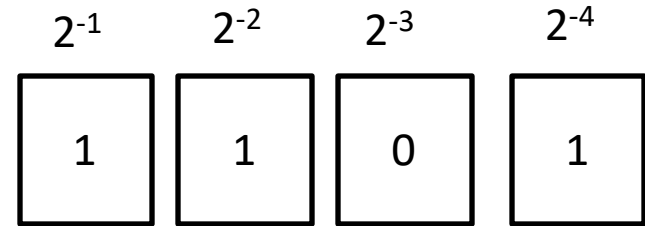


CONVERSION EXAMPLE

Let's convert 0.8125 to floating-point representation

$$\begin{array}{l} 0.8125 \times 2 = 1.6250 \quad \mathbf{1} \\ 0.6250 \times 2 = 1.2500 \quad \mathbf{1} \\ 0.2500 \times 2 = 0.5000 \quad \mathbf{0} \\ 0.5000 \times 2 = 1.0000 \quad \mathbf{1} \end{array}$$


$$\begin{array}{l} 0.1101 \\ = 1.101 \times 2^{-1} \end{array}$$



$$1 \times 1/2 + 1 \times 1/4 + 0 \times 1/8 + 1 \times 1/16 = 13/16$$

CONVERSION EXAMPLE PART 2

$$0.1101 = 1. \boxed{101} \times \boxed{2^{-1}}$$

Sign: 0

Mantissa: 101

Exponent: $-1 + 127 = 126(d)$
 $= 01111110(b)$

0 01111110 1010000 00000000 00000000

CONVERSION

0.1 x 2 = 0.2	0
0.2 x 2 = 0.4	0
0.4 x 2 = 0.8	0
0.8 x 2 = 1.6	1
0.6 x 2 = 1.2	1
0.2 x 2 = 0.4	0

..... repeats ...

Just like the 1/3, 0.1 keeps repeating in binary

0 01111011 1001100 11001100 1100110

$$123 \quad \frac{1}{2} + \frac{1}{2^4} + \frac{1}{2^5} + \frac{1}{2^8} + \frac{1}{2^9} + \frac{1}{2^{12}} + \frac{1}{2^{13}} + \frac{1}{2^{16}} + \frac{1}{2^{17}} + \frac{1}{2^{20}} + \frac{1}{2^{21}}$$

$$(-1)^s \times (1 + m) \times 2^{\text{exponent} - \text{bias}}$$

$$0.09999999940395355224609375 = (-1)^0 \times \left(1 + \frac{1258291}{2097152}\right) \times 2^{123 - 127}$$

Not quite 0.1

EXAM QUESTION

CS 2130 Exam 1, Spring 2023

Page 5 of 6

UVA computing id:

Page 5: Floating Point and Bitwise Operations

8. [9 points] Write the following binary number as an 8-bit floating point number assuming a 4-bit exponent value.

+0.00001110011

Answer

EXAM QUESTION

CS 2130 Exam 1, Spring 2023

Page 5 of 6

UVA computing id:

Page 5: Floating Point and Bitwise Operations

8. [9 points] Write the following binary number as an 8-bit floating point number assuming a 4-bit exponent value.

+0.00001110011

Answer

0 0010 110

TODAY'S LECTURE

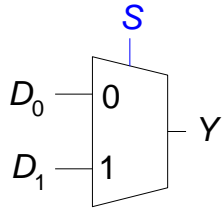


Contents

1. Introduction to Muxes
2. Introduction to Clocks.

MUXES

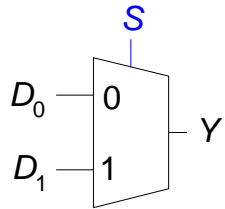
Example: 2:1 Mux



S	D_1	D_0	Y	S	Y
0	0	0	0	0	D_0
0	0	1	1	1	D_1
0	1	0	0		
0	1	1	1		
1	0	0	0		
1	0	1	0		
1	1	0	1		
1	1	1	1		

- Selects between one of N inputs to connect to output
- **Select** input is $\log_2 N$ bits – control input

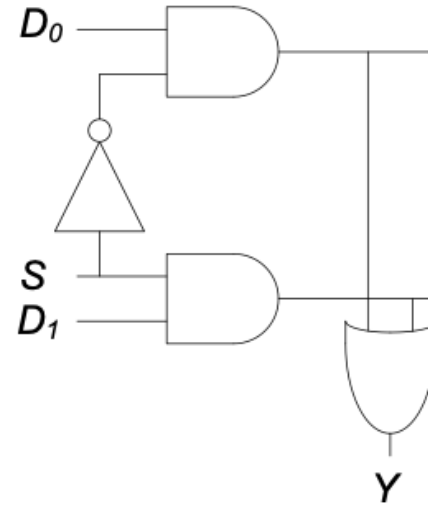
1 BIT MUX



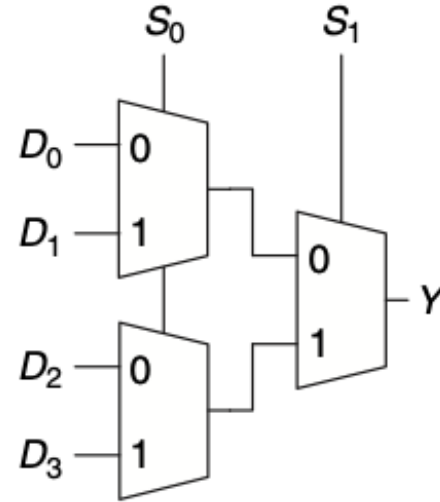
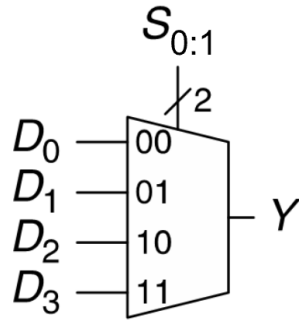
$$Y = D_0\bar{S} + D_1S$$

S	D ₁	D ₀	Y
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

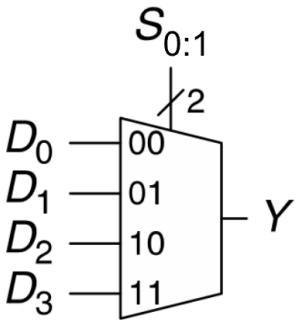
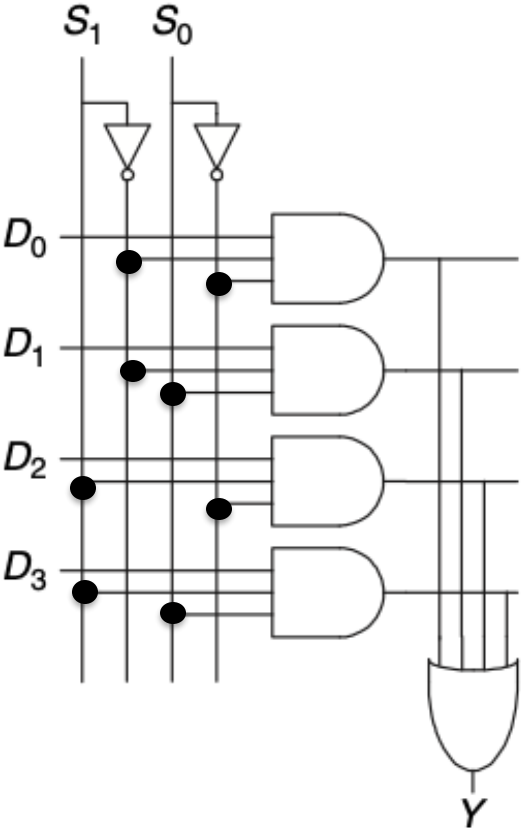
S	Y
0	D ₀
1	D ₁



2 BIT MUX



2 BIT MUX

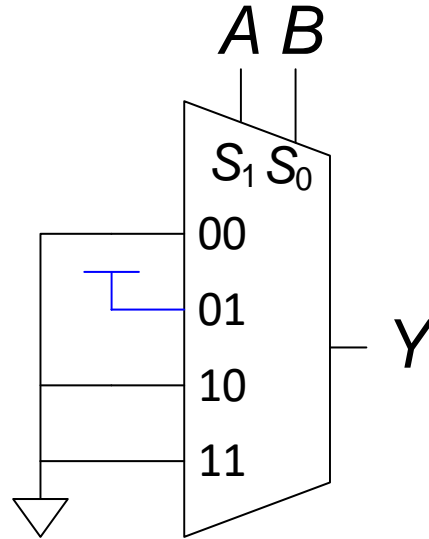


MUX AS A LOOK UP TABLE

Using mux as a **lookup table**

A	B	Y
0	0	0
0	1	1
1	0	0
1	1	0

$$Y = AB$$

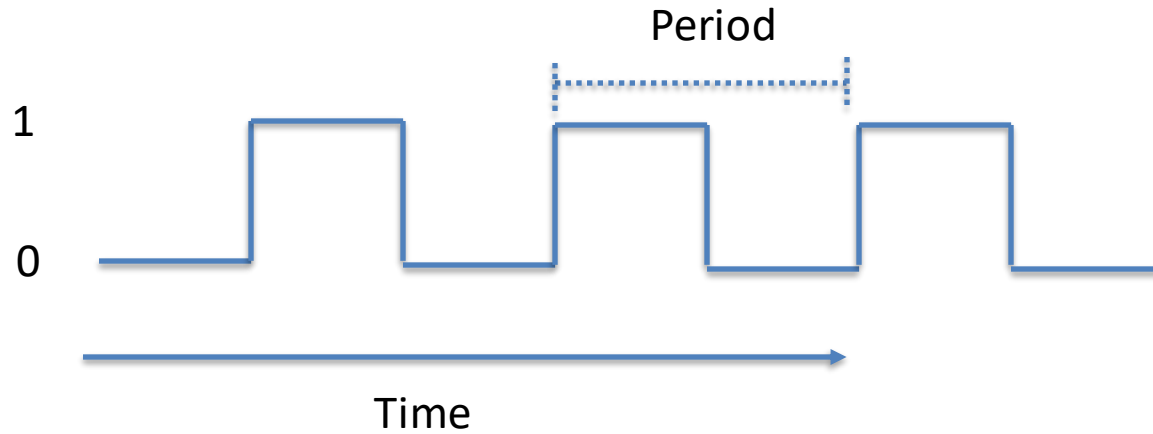


CLOCKS

A clock is something that produces a periodic signal

Period is length of time for one clock cycle

Example

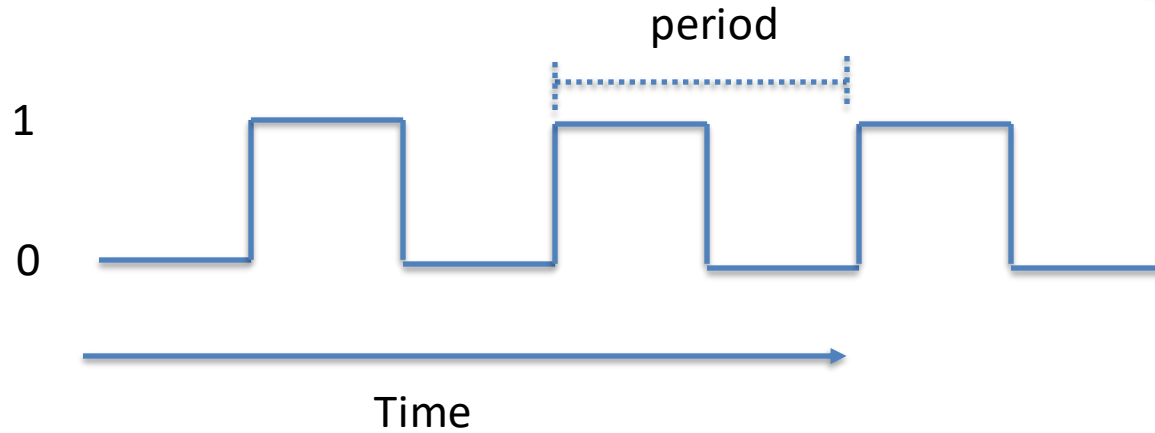


CLOCKS

Clock frequency Intel Core i-9 3.0GHz

Frequency = $1/\text{period}$

Period = $1/\text{frequency}$

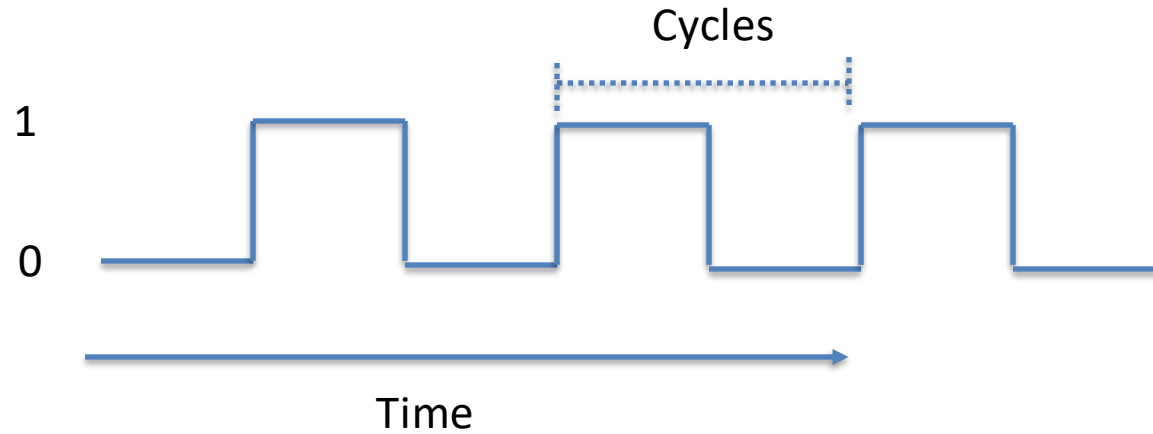


CLOCKS

Clock frequency Intel Core i-9

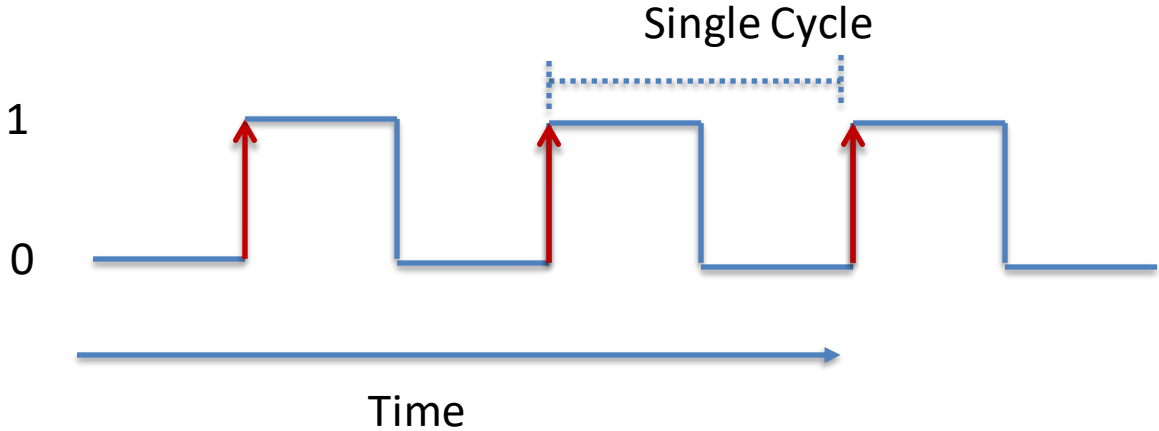
3.0GHz. = $3 \times 10^9 = 3,000,000,000$ cycles per second

That is fast.



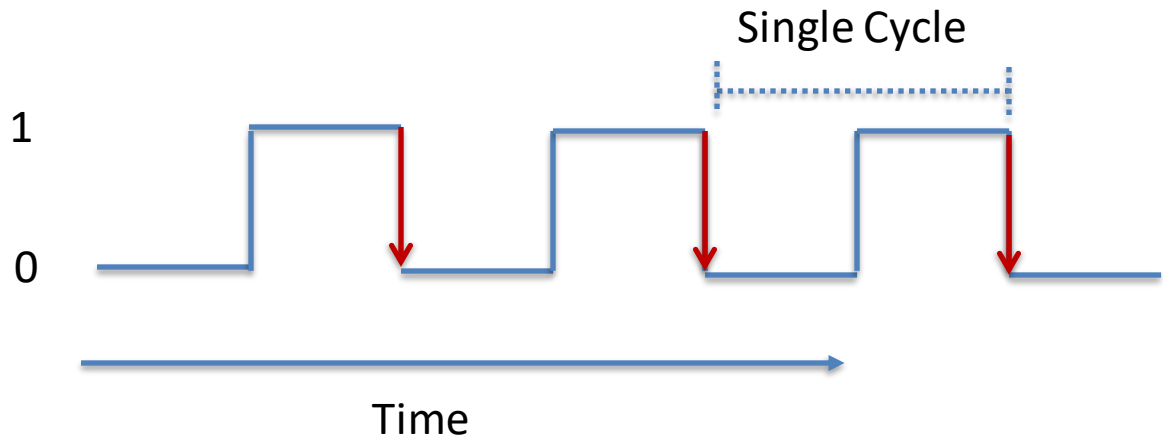
CLOCK EDGES

Rising Edge



CLOCK EDGES

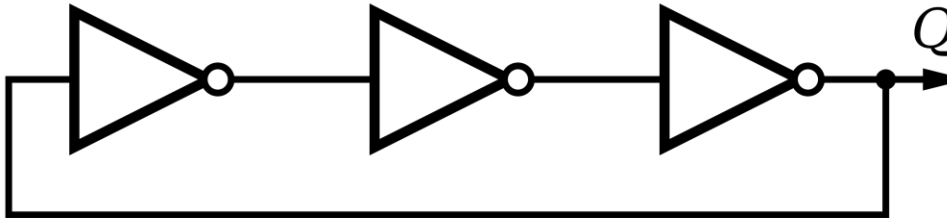
Falling Edge.



We will build a single cycle machine. It will complete all the computation in a single cycle.

USING RING OSCILLATORS TO GENERATE CLOCKS

A clock is something that produces a periodic signal



Let's walk through
an example
assume that Q
starts off as 0

$$\text{Frequency} = 1/(2*t*n)$$

Where t is the time delay of an inverter and
n is number of inverters

